Optimization Project in Energy ENT306

Elise Grosjean Ensta-Paris

Ensta-Paris Institut Polytechnique de Paris



▲ロ ▶ ▲周 ▶ ▲ 国 ▶ ▲ 国 ▶ ● の Q @

1 Introduction

2 Deterministic model



Introduction

 Main goal: programming numerical methods (Energy Management System) for a very simple model of a microgrid.

• Microgrid: a set consisting of the following elements:

- a small-size electrical load (a building, a few houses)
- a source of renewable energy (solar panels, a wind turbine)
- a storage device (Energy Storage System, a battery)
- a macrogrid (an access to a large-scale energy network).

Introduction

- Time scale: the decisions are taken every hour during the day.
- Optimization variables (at each time step) :
 - the amount of energy to be stored or withdrawn from the battery
 - the amount of energy **bought** or **sold** to the network.

Contraints:

- nonnegativity of variables
- evolution of the battery
- storage capacity of the battery.

Cost function:

- Cost of the energy bought on the network...
- minus the cost of the energy sold on the network.

Introduction

Main challenge:

- The electrical load and the production of renewable energy are random.
- No available probabilistic model, instead the history of electrical load and solar production.

Tools and mathematical concepts:

- Dynamic programming (temporal aspect for decision process).
- Autoregressive processes.

Optimization is useful for...

- minimizing the management costs
- modelling random processes
- describing some functions, for which no analytical expression is available (interpolation).

Introduction

Philosophy:

- **Compromise** between the complexity of stochastic modelling and solvability of the problem.
- Emphasis on the mathematical approach. Maths concepts useful in other application contexts.
- We will work with models of **increasing** complexity.

Warning:

- (Very) simplified models, with artificial data.
- The purpose is not to conclude on the relevance of the introduction of such or such technology.

Introduction

References :

- Le Franc, Carpentier, Chancelier, de Lara. EMSx: an Energy Management System numerical benchmark. Energy System, 2021.
- Hafiz, Awal, de Queiroz, Husain. Real-time Stochastic Optimization of Energy Storage Management using Rolling Horizon Forecasts for Residential PV Applications, 2019.
- Olivares et al. Trends in microgrid control. IEEE Transactions on smart grid, 2014.

▲ロ ▶ ▲周 ▶ ▲ 国 ▶ ▲ 国 ▶ ● の Q @

Introduction

Organisation:

- 6 units: January 17, January 07, January 14, January 21, January 28, February 4.
- Work in pairs (please form the groups by next week).
- Programming with Python.
- Evaluation: programming exercises to solve + work in class.

1 Introduction

2 Deterministic model



- Horizon: 24 hours, stepsize: 1 hour. Optimization over T = 24 intervals.
- Optimisation variable :
 - x(s) : state of charge of the battery at time s, s = 1, ..., T + 1
 - *a*(*s*): amount of electricity bought on the network (*s* = 1, ..., *T*).
 - v(s): amount of energy sold on the network (s = 1, ..., T).

Parameters:

- d(s): net demand of energy (load minus solar production) at time s, s = 1, ..., T.
- $P_a(s)$: unitary buying price of energy at time s
- $P_v(s)$: unitary selling price of energy at time s
- *x*_{max}: storage capacity of the battery.

Remark: the demand is supposed to be deterministic (that is to say, known in advance), for the moment.

Contraints:

$$\begin{array}{l} \mathbf{x}(s+1) = x(s) - d(s) + a(s) - v(s), \ \forall s = 1, ..., T \\ \mathbf{x}(1) = 0 \\ \mathbf{a}(s) \geq 0, \ \forall s = 1, ..., T \\ \mathbf{v}(s) \geq 0, \ \forall s = 1, ..., T \\ \mathbf{0} \leq x(s) \leq x_{\max}, \ \forall s = 1, ..., T + 1. \end{array}$$

Cost function to be minimized:

$$J(x,a,v) = \sum_{s=1}^{T} \left(P_a(s)a(s) - P_v(s)v(s) \right).$$

The buying and selling prices P_a and P_v depend on time. It holds: $P_a(s) > P_v(s)$, so that it is useless to try to buy and sell electricity on the network at the same time!

▲ロ ▶ ▲周 ▶ ▲ 国 ▶ ▲ 国 ▶ ● の Q @

Deterministic model

Exercise 1

Write and solve the optimization problem in a form that is compatible with the function linprog of Python.

Main idea behind dynamic programming:

- We **parametrize** the problem to be solved ~→ a sequence of problems of increasing complexity.
- We look for a relation ("dynamic programming principle") between the optimal values of the different problems.

Parameters :

- Initial time $t \in \{1, ..., T + 1\}$.
- Initial state-of-charge of the battery $y \in [0, x_{max}]$.

We are interested in the problem with t = 1 and y = 0.

うしん 同一人用 イモット 一切 くう

Parameterized problem:

$$V(t, y) = \inf_{\substack{x(t), x(t+1), \dots, x(T+1) \\ a(t), a(t+1), \dots, a(T) \\ v(t), v(t+1), \dots, v(T)}} \sum_{s=t}^{T} P_{a}(s)a(s) - P_{v}(s)v(s)$$
(P(t, y))

under the constraints:

•
$$x(s+1) = x(s) - d(s) + a(s) - v(s), \forall s = t, ..., T$$

• $x(t) = y$
• $a(s) \ge 0, \forall s = t, ..., T$
• $v(s) \ge 0, \forall s = t, ..., T$
• $0 \le x(s) \le x_{max}, \forall s = t, ..., T + 1.$

The function V is called **value function**; it plays a crucial role, in particular in the treatment of the stochastic version of the problem.

Parameterized problem:

$$V(t, y) = \inf_{\substack{x(t), x(t+1), \dots, x(T+1) \\ a(t), a(t+1), \dots, a(T) \\ v(t), v(t+1), \dots, v(T)}} \sum_{s=t}^{T} P_{a}(s)a(s) - P_{v}(s)v(s)$$
(P(t, y))

under the constraints:

•
$$x(s+1) = x(s) - d(s) + a(s) - v(s), \forall s = t, ..., T$$

• $x(t) = y$
• $a(s) \ge 0, \forall s = t, ..., T$
• $v(s) \ge 0, \forall s = t, ..., T$
• $0 \le x(s) \le x_{max}, \forall s = t, ..., T + 1.$

The function V is called **value function**; it plays a crucial role, in particular in the treatment of the stochastic version of the problem.

- Dynamic programming = Temporal decomposition with smaller sub-problems
- Dynamic programming principle = Recursive relation
 - process ensuring that each subproblem is solved optimally, taking into account future implications and not the previous decisionst.
- Interests of the decomposition:
 - reducing complexity
 - optimality of the global solution thanks to Bellman optimality principle
 - clearer implementation
 -

for a

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 のへの

Deterministic model and dynamic programming

Theorem [Dynamic programming principle]

The following holds true:

$$V(t, y) = \inf_{\substack{(z, a, v) \in \mathbb{R}^3 \\ (z, a, v) \in \mathbb{R}^3}} P_a(t)a - P_v(t)v + V(t+1, z),$$

$$\begin{cases} a \ge 0 \\ v \ge 0 \\ z = y - d(t) + a - v \\ 0 \le z \le x_{\max} \end{cases}$$
(DP(t, y))
all $t \in \{1, ..., T\}$ and $y \in [0, x_{\max}]$ and
$$V(T+1, y) = 0.$$

$$V(t,y) = \inf_{(z,a,v) \in \mathbb{R}^3} P_a(t)a - P_v(t)v + \frac{V(t+1,z)}{V(t+1,z)}$$



$$V(t,y) = \inf_{(z,a,v)\in\mathbb{R}^3} P_a(t)a - P_v(t)v + V(t+1,z)$$



$$V(t,y) = \inf_{(z,a,v)\in\mathbb{R}^3} P_a(t)a - P_v(t)v + \frac{V(t+1,z)}{V(t+1,z)}$$





Deterministic model

Computation of the value function

- Let us suppose that the function $z \mapsto V(t+1, z)$ is known, with $t \in \{1, ..., T\}$. By solving DP(t, y) for all possible values of y, we can compute $y \mapsto V(t, y)$.
- In practice, we can only solve DP(t, y) for a sample of values of y ∈ [0, x_{max}]. We evaluate V(t, y) for those values of y; then we interpolate.
- In that way, we can compute (an approximation of) the value function, in a recursive and backward fashion (from T + 1 to 1).

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ のQで

Deterministic model

Solving the initial problem.

- Let (z, a, v) denote a solution of DP(1, 0). Let us set: $\bar{x}(2) = z$, $\bar{a}(1) = a$, $\bar{v}(1) = v$.
- Let (z, a, v) be a solution of $DP(2, \overline{x}(2))$. Let us set: $\overline{x}(3) = z$, $\overline{a}(2) = a$, $\overline{v}(2) = v$.
- Let (z, a, v) be a solution of $DP(3, \overline{x}(3))$. Let us set: $\overline{x}(4) = z$, $\overline{a}(3) = a$, $\overline{v}(3) = v$.
- ... and so on, until the resolution of $DP(T, \bar{x}(T))$.

0

- Let $y \in \mathbb{R}^J$ and let $val = V(t, y) \in \mathbb{R}^J$. We consider a function f such that $val_j = f(y_j)$, for all j = 1, ..., J.
- We interpolate f with a second-order polynomial, by solving

$$\inf_{\alpha \in \mathbb{R}^3} \sum_{j=1}^J (\alpha_1 + \alpha_2 y_j + \alpha_3 y_j^2 - \mathsf{val}_j)^2.$$

• Let $\bar{\alpha}$ be the solution. We get the approximation:

$$f(y) \approx \bar{\alpha}_1 + \bar{\alpha}_2 y + \bar{\alpha}_3 y^2.$$

Exercise 2

Write a function interpolate implementing the interpolation with a second-order polynomial. Inputs: J, $y \in \mathbb{R}^J$, $val_j \in \mathbb{R}^J$. Outputs: $\alpha \in \mathbb{R}^3$.

Deterministic model

Exercise 3

Write a function DP_solve with output the solution (z, a, v) of problem DP(t, y).

We assume that the function $V(t + 1, \cdot)$ is approximated by a second-order polynomial, described by a coefficient $\alpha \in \mathbb{R}^3$.

Input: $t \in \{1, ..., T\}$, $y \in [0, x_{\max}]$, $\alpha \in \mathbb{R}^3$.

Deterministic model

Exercise 4

Write a function DP_backward which compute a polynomial approximation of V, in the form of a matrix $\alpha \in \mathbb{R}^{(T+1)\times 3}$, that is to say:

$$V(t,y) \approx \alpha_{t,1} + \alpha_{t,2}y + \alpha_{t,3}y^2.$$

We will proceed in a recursive fashion:

- Given $\alpha_{t+1,1}$, $\alpha_{t+1,2}$, $\alpha_{t+1,3}$, evaluate $V(t, y_j)$ for all j = 1, ..., J, where $y_j = (j-1)/(J-1)x_{max}$.
- Calculate $\alpha_{t,1}$, $\alpha_{t,2}$, and α_{t_3} by interpolating the values of $V(t, \cdot)$.

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 のへの

Deterministic model

Exercise 5

Write a function DP_forward computing the solution to the original problem. We will first make use of DP_backward to get an approximation of the value function.